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# CS6233 Midterm

Answer the questions in the spaces provided in the question sheets. If you run out of room for an answer, you can continue on the back of the page.

|  |  |  |
| --- | --- | --- |
| Question | Points | Score |
| 1 | 10 |  |
| 2 | 10 |  |
| 3 | 10 |  |
| 4 | 10 |  |
| 5 | 10 |  |
| 6 | 10 |  |
| 7 | 16 |  |
| 8 | 16 |  |
| 9 | 14 |  |
| Total | 106 |  |

1. Consider the following C program:

#include <stdio.h>

#include <stdlib.h>

#include <unistd.h>

#include <sys/wait.h>

int main**(**void**)** **{** int i **=** 0**;** **for** **(**i **=** 0**;** i **<** 4**;** i**++)** **{** **if** **(**fork**()** **==** 0**)** **{** printf**(**"foo\n"**);**

## **}** **else** **{** wait**(NULL);** exit**(**1**);**

**}**

**}**

**return** 0**;**

**}**

1. (5 points) How many times will foo" be printed?

1. (5 points) How many processes will be created (including the initial process)?

2. The following is the code that implements the kill system call in xv6:

Assembly Code:

1 kill **:**

## 2 mov **$**0x6 ,% eax 3 int **$**0x40 4 ret

C code:

1. // Fetch the int at addr from the current process .
2. int
3. fetchint **(** uint addr **,** int **\***ip**)**
4. **{**

## 5 **if(** addr **>=** proc **->**sz **||** addr **+**4 **>** proc **->**sz**)**

1. **return** **-**1**;**
2. **\***ip **=** **\*(** int **\*)(** addr **);**
3. **return** 0**;**
4. **}**

10

1. // Fetch the nth 32- bit system call argument .
2. int
3. argint **(**int n**,** int **\***ip**)**
4. **{**

## 15 **return** fetchint **(**proc **->**tf **->**esp **+** 4 **+** 4**\***n**,** ip **);**

16 **}**

17

1. int
2. sys\_kill **(** void **)**
3. **{**
4. int pid **;**

22

1. **if(** argint **(**0**,** **&**pid **)** **<** 0**)**
2. **return** **-**1**;**

## 25 **return** kill **(**pid **);**

26 **}**

1. (3 points) On line 15 of the C code, in the argint function, why do we add 4 to proc->tf->esp?

1. (2 points) On that same line, why do we multiply n by 4?

1. (5 points) What could go wrong if xv6 didn't have the check at line 5 in fetchint?

3. Three batch jobs, A through C, arrive at a computer center at the same time. They have estimated running times of 10, 6, 2 minutes respectively. At time 3, jobs D and E arrive, which take 4 and 8 minutes respectively. For each of the following scheduling algorithms, determine the mean process turnaround time. Ignore process switching overhead.

1. (5 points) First-come, first-served (run in order A, B, C, D, E).
2. (5 points) Shortest job first.

4. Recall that in 32-bit x86, page directories and page tables are each made up of 1024 32-bit entries. Suppose we have 4 processes on a system, each of which has every possible virtual address mapped.

1. (5 points) How much memory is used to store the page directories and page tablesif 4KB pages are used?
2. (5 points) If 4MB pages (super pages) are used, then the entries in the page direc- tory point directly to the page frame (i.e., no second-level page tables are used). How much memory would be taken up by page directories in this case?

5. Below is the code for swtch() from xv6:

## 1 # Context switch 2 # 3 # void swtch ( struct context \*\* old , struct context \*new); 4 #

5 # Save current register context in old 6 # and then load register context from new.

7

1. .globl swtch
2. swtch **:**
3. movl 4**(**% esp ), %eax 11 movl 8**(**% esp ), %edx

12

## 13 # Save old callee - save registers 14 pushl %ebp 15 pushl %ebx 16 pushl %esi 17 pushl %edi

18

19 # Switch stacks

## 20 movl %esp , (% eax **)** 21 movl %edx , %esp

22

23 # Load new callee - save registers

## 24 popl %edi 25 popl %esi 26 popl %ebx 27 popl %ebp 28 ret

1. (5 points) How is the program counter (%eip) saved and restored by this code?

1. (5 points) Give an example of one other piece of program state that is not explicitly saved by this code and explain why that doesn't cause problems when performing a context switch.

6. On the next page is the code for the scheduler function in xv6, which picks the next process to run. Refer to it as you answer the following questions:

1. (4 points) What scheduling algorithm is implemented by this code? Is it fair (i.e., does each process get an equal share of the CPU)?

1. (4 points) On lines 20 and 21, we skip the rest of the loop if the process state is not RUNNABLE. What is an example of another state that a process could be in? What would go wrong if we tried to run a process in that state?

1. (2 points) The swtch function, called on line 29, is written in assembly. What is its purpose, and why does it need to be written in assembly rather than C?

1. // Per -CPU process scheduler .
2. // Each CPU calls scheduler () after setting itself up.
3. // Scheduler never returns . It loops , doing :
4. // - choose a process to run
5. // - swtch to start running that process
6. // - eventually that process transfers control
7. // via swtch back to the scheduler .
8. void

## 9 scheduler **(** void **)**

10**{**

11 struct proc **\***p**;**

12

1. **for** **(;;){**
2. // Enable interrupts on this processor .

## 15 sti **();**

16

17 // Loop over process table looking for process to run .

## 18 acquire **(&** ptable **.** lock **);** 19 **for** **(**p **=** ptable **.** proc **;** p **<** **&** ptable **.** proc **[** NPROC **];** p **++){** 20 **if(**p**->** state **!=** RUNNABLE **)**

21 **continue** **;**

22

23 // Switch to chosen process . It is the process 's job 24 // to release ptable . lock and then reacquire it

25 // before jumping back to us.

## 26 proc **=** p**;** 27 switchuvm **(**p**);** 28 p**->** state **=** RUNNING **;** 29 swtch **(&** cpu **->** scheduler **,** proc **->** context **);** 30 switchkvm **();**

31

1. // Process is done running for now .
2. // It should have changed its p-> state before coming back .
3. proc **=** 0**;**
4. **}**

## 36 release **(&** ptable **.** lock **);**

37

1. **}**
2. **}**

7. (16 points) For the following questions Circle True/False. Each question is worth 2 points

T F (a)  A process that is not in main memory is immediately available for execution, regardless

of whether or not it is awaiting an event.

T F (b)  For virtual memory, a hardware mechanism is needed for translating relative addresses

to physical main memory addresses at the time of execution of the instruction that contains

the reference.

T F (c)  The size of virtual storage is limited by the actual number of main storage locations.

T F (d)  It is the responsibility of the operating system to control the execution of processes.

T F (e)  The OS may suspend a process if it detects or suspects a problem.

T F (f)  If a system does not employ virtual memory each process to be executed must be fully

loaded into main memory.

T F (g)  An example of an application that could make use of threads is a browser.

T F (h)  Segmentation is not visible to the programmer.

1. (16 points) For the following questions circle the correct option for the answer. Each question is worth 2 points

* 1. Main memory divided into a number of equal size frames is the \_\_\_\_\_\_\_\_\_\_ technique.
     1. simple paging   B)  dynamic partitioning

C)  fixed partitioning   D)  virtual memory segmentation

* 1. The chunks of a process are known as \_\_\_\_\_\_\_\_\_\_ .
     1. pages   B)  addresses

C)  frames   D)  segments

* 1. Available chunks of memory are known as \_\_\_\_\_\_\_\_\_\_\_ .
     1. frames   B)  segments

C)  addresses   D)  pages

* 1. The four main structural elements of a computer system are:
     1. Processor, Main Memory, I/O Modules and System Bus
     2. Processor, I/O Modules, System Bus and Secondary Memory
     3. Processor, Registers, Main Memory and System Bus
     4. Processor, Registers, I/O Modules and Main Memory

* 1. The \_\_\_\_\_\_\_\_\_\_ holds the address of the next instruction to be fetched.
     1. Accumulator (AC)   B)  Instruction Register (IR)

C)  Instruction Counter (IC)   D)  Program Counter (PC)

* 1. Instruction processing consists of two steps:
     1. fetch and execute   B)  instruction and execute

C)  instruction and halt   D)  fetch and instruction

* 1. The processor itself provides only limited support for multiprogramming, and \_\_\_\_\_\_\_\_\_\_ is needed to manage the sharing of the processor and other resources by multiple applications at the same time.

* + 1. memory   B)  data

C)  software   D)  hardware

* 1. When one process spawns another, the spawned process is referred to as the \_\_\_\_\_\_\_\_\_\_ .

* + 1. zombie process   B)  child process

C)  stack process   D)  parent process

1. (16 points) For the following questions fill in the blank with the correct answer. Each question is worth 2 points.

* 1. \_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_ is a storage allocation scheme in which secondary memory can be addressed as though it were part of main memory.

* 1. To overcome the problem of doubling the memory access time, most virtual memory schemes make use of a special high-speed cache for page table entries called a \_\_\_\_\_\_\_\_\_\_ .

* 1. A process in the \_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_ state is in main memory and available for execution.

* 1. When the system spends most of its time swapping pieces rather than executing instructions it leads to a condition known as \_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_ .

* 1. When instead of using one page table entry per virtual page, the system keeps one entry per physical page frame, i called: \_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_

* 1. My **favorite** part of the course so far has been:

\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_ \_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_

* 1. My **least**-favorite part of the course so far has been (besides this test):

\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_

\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_\_